

Tutorial #11
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Introduction to Data Communications

1. Obtain the Channel Capacity in bit rate, on a 2500 Hz carrier, if the signal to noise ratio (SNR)dB = 20 dB.

Solution:

Given,

$$(SNR) \text{ dB} = 20\text{dB}$$

Using the formula,

$$(SNR) \text{ dB} = 10 * \log_{10} (\text{signal power/noise power})$$

Therefore,

$$20\text{dB} = 10 * \log_{10} (snr)$$

$$\log_{10} (snr) = 20/10$$

$$= 2$$

$$snr = 10^2$$

$$= 100$$

Using Shannon's Law,

$$C = B * \log_2 [1 + snr]$$

$$C = 2500 * \log_2 (1 + 100)$$

$$C = 2500 * \log_2 (101) \quad ; \log_2 (101) = \frac{\log_{10}(101)}{\log_{10}(2)}$$

$$= 2500 * 6.66$$

$$C = 16.65 \text{ Kbps}$$

2. Obtain the CRC for the data stream 0100 0011 ('C') using pattern 1011 using the shifted polynomial method.

Solution:

(Data polynomial form * highest power of CRC ($=x^3$))

a) First shift the input data by 3 bits to make room for 3 CRC bits (the CRC pattern 1011, has 4 bits hence CRC is $4-1 = 3$ bits)

(Another way of looking highest power in CRC pattern is x^3 - hence CRC has 3 bits)

Data **0 1 0 0 0 1 1**

Shifted data (by 3) 0 1 0 0 0 1 1 0 0 0

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CRC bits init to 0

Shifted data $D(x) = x^9 + x^4 + x^3$

****** a simple method, when they are more experienced, shift left by 3 = multiply by x^3**

Shifted $D'(x) = (x^6 + x + 1) * x^3 = x^9 + x^4 + x^3$

b) then calculate CRC for the shifted data polynomial

$$\begin{array}{r}
 x^6 + x^4 + x^3 + x^2 + x \\
 x^3 + x + 1 \overline{) x^9 + x^7 + x^6 + x^4 + x^3} \\
 \underline{x^9 + x^7 + x^6} \\
 x^7 + x^6 + x^4 + x^3 \\
 \underline{x^7 + x^5 + x^4} \\
 x^6 + x^5 + x^3 \\
 \underline{x^6 + x^4 + x^3} \\
 x^5 + x^4 \\
 \underline{x^5 + x^3 + x^2} \\
 x^4 + x^3 + x^2 \\
 \underline{x^4 + x^2 + x} \\
 x^3 + x \\
 \underline{x^3 + x + 1} \\
 1
 \end{array}$$

$$CRC(x) = 1$$

$$\Rightarrow CRC \text{ bits} = 001 \text{ (3-bit CRC)}$$

c) Then insert the CRC in the CRC bits created by the shift (a)

Sent message bits:

Initially:

Shifted data (by 3) 01000011000

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CRC bits init to 0

We replace the CRC bits (000) by calculated CRC 001

Shifted data (by 3) 01000011001

|----|

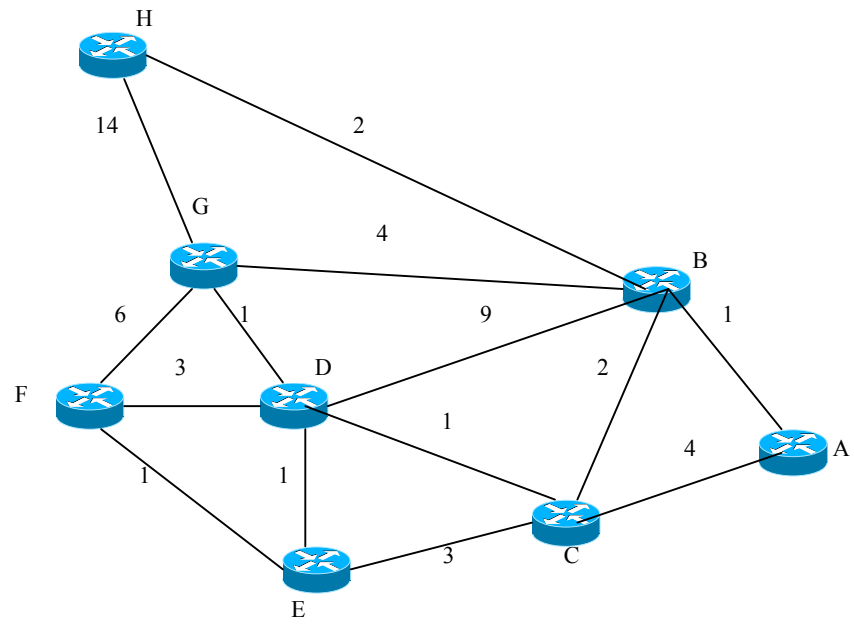
CRC bits as calculated

Transmitted message = 01000011001

(Simple method)

$$Sent \text{ message} = D(x) * x^3 + CRC(x) = x^9 + x^4 + x^3 + 1$$

3. Using Dijkstra's algorithm, compute the shortest path from node E.



Solution: Dijkstra's at Node E

$\{M\}$	$L(A)$	$Path$	$L(B)$	$Path$	$L(C)$	$Path$	$L(D)$	$Path$
$\{E\}$	∞	----	∞	----	3	E-C	1	E-D
$\{E,D\}$	∞	----	10	E-D-B	2	E-D-C	1	E-D
$\{E,D,F\}$	∞	----	10	E-D-B	2	E-D-C	1	E-D
$\{E,D,F,C\}$	6	E-D-C-A	4	E-D-C-B	2	E-D-C	1	E-D
$\{E,D,F,C,G\}$	6	E-D-C-A	4	E-D-C-B	2	E-D-C	1	E-D
$\{E,D,F,C,,G,B\}$	5	E-D-C-B-A	4	E-D-C-B	2	E-D-C	1	E-D
$\{E,D,F,C,G,B,A\}$	5	E-D-C-B-A	4	E-D-C-B	2	E-D-C	1	E-D
$\{E,D,F,C,G,B,A,H\}$	5	E-D-C-B-A	4	E-D-C-B	2	E-D-C	1	E-D
$\{E,D,F,C,G,B,A,H\}$	5	E-D-C-B-A	4	E-D-C-B	2	E-D-C	1	E-D

$L(F)$	$Path$	$L(G)$	$Path$	$L(H)$	$Path$
1	E-F	∞	---	∞	---
1	E-F	2	E-D-G	∞	---
1	E-F	2	E-D-G	∞	---
1	E-F	2	E-D-G	∞	---
1	E-F	2	E-D-G	16	E-D-G-H
1	E-F	2	E-D-G	6	E-D-C-B-H
1	E-F	2	E-D-G	6	E-D-C-B-H
1	E-F	2	E-D-G	6	E-D-C-B-H
1	E-F	2	E-D-G	6	E-D-C-B-H